

# Bewl

**a programming language for topos theory**

more precisely, a Scala DSL

for the Mitchell-Benabou internal language of a topos,  
with some topos implementations

# Summary

- About me
- Topos theory: a promising conceptual tool for escaping the limitations of set-based math
- It needs a "four function calculator":  
this is why I wrote Bewl
- Engineering compromises:  
Topos theory on a finite machine
- What can Bewl do so far?
- Future directions

# About me, Felix Dilke

- I'm not an academic
- I'm a software developer on SpringerLink ([link.springer.com](http://link.springer.com))
- Bewl is my 10% time project at Springer Nature
- I've presented it to my colleagues, who are experienced real-world software developers with an interest in science
- So, this talk will be informal, even impressionistic, with many speculative analogies between software and math, and I apologize in advance to domain experts if I seem to be making wild claims.
- Bewl is test-driven Scala code, open source on GitHub

# The limitations of set-based math

A whole talk in itself, but briefly:

Colin McLarty's book on topoi has led me to see set theory as a legacy platform like MS-DOS (!) with many limitations and anomalies which topos theory can explain and perhaps alleviate. Examples:

- Large cardinals: a theory peculiar to sets, with unclear application to mainstream math.
- Ultrafilters: chimerical objects
- Permutation parity: an unexplained feature of the topos of sets (another motivation for Bewl was to explain parity).

# Topos theory

- A topos is a category with all the optional extras - finite products, equalizers, exponents, subobject classifier
- Inside a topos, one can define algebraic structures, quantifiers and a near-classical internal logic. So the topos becomes a workspace in which one can do math.
- A topos is an abstraction of the category of sets (the usual foundation for math) and shows how to do much of the same work in a much wider context.
- Driving metaphor: a topos is a "virtual machine, for math"
- Definitions, constructions, theorems "run" in a topos just as apps run on a VM, or SQL statements run on a database

# The promise

- Cleanly separate language from implementation  
(just as webdesigners separate HTML from business logic)
- A lot of math can easily be "refactored" to apply  
in a much wider context
- Example: Schur's lemma (that the endomorphisms of a  
simple module form a division ring) can be expressed  
in topos-valid form
- Understand and escape limitations of set-based reasoning

# Examples of topoi

- Sets. Also finite sets (foundational for Bewl)
- Smooth sets (a workspace for synthetic differential geometry)
- The effective topos (a workspace for computability)
- Sheaves (workspaces for algebraic geometry)
- Fuzzy sets (with equality taking values in a Heyting algebra)
- Diagrams of a given shape:

graphs

automorphisms

monoid actions

# Example: Schur's lemma

In an arbitrary topos, it's still true that endomorphisms of a simple module form a division ring

- Define a ring  $R$  as an object with arrows  $*$ :  $R \times R \rightarrow R$ , unit:  $1 \rightarrow R$ , subject to certain laws as arrow equalities
- Similarly an  $R$ -module has an abelian group structure and conditioned scalar multiplication  $**$ :  $R \times M \rightarrow M$
- "M is simple" is then expressed using a quantifier over  $M$ : all submodules of  $M$ , i.e. subobjects of  $M$  obeying certain closure laws, are equal to either  $0$  or  $M$
- One can then formally construct the ring of endomorphisms as a subobject of the exponential object  $M \wedge M$ , and show it obeys a law "for all  $x$ , either  $x = 0$  or  $x$  has an inverse" in topos terms, making it a division ring.

## Example: Schur's lemma (2)

The details of all this involve the Mitchell-Benabou internal language, which interprets logical formulas as statements about equality between arrows in the topos. A soundness result formalizes the proof as pure symbol-manipulation.

So now we have Schur's lemma for graphs, sheaves, fuzzy sets, monoid actions, etc.

We've also separated the language (abstract strings of symbols) from the implementation (specific topoi).

Arguably, this yields a cleaner and more definitive version of the original result.

# Refactoring math

Large blocks of math translate similarly without much change.

As a more elaborate example, one can do the same for the Los ultraproduct theorem (Volger 1975).

Once can also do topology in a topos, using Manes' theorem (the algebra of compact Hausdorff spaces) as a starting point.

This points to an ambitious, Hilbertian program to refactor math - realized, I believe, in homotopy type theory (HoTT)

Bewl is just a DSL ("four function calculator"), but already many of its library methods are software versions of definitions and constructions from topos theory.

# Aggressively refactoring the foundations of math:

Obstacles:

- Topos logic is *intuitionistic*, i.e. allows multiple truth-values and no excluded middle
- Much classical math is irretrievably Boolean and can't easily be rewritten this way  
Example: number theory
- Basic concepts like finiteness and the real numbers don't have immediately clear analogues

There are potential answers to these, and they involve fascinating conceptual questions (see HoTT).

- But most of all: How do you do the calculations?

## Example: music theory

"The Topos of Triads", Thomas Noll, 2005

(paper on CiteSeer)

"The Topos of Music", Guerino Mazzola, 2002

(a book on SpringerLink)

Noll explores music theory by defining a "triadic monoid" and working in the topos of actions over it.

He had to do all these calculations by hand, for example enumerating topologies on the triadic topos.

Bowl can now do many of these computations itself.

In particular, I verified that the C major chord

(modelled as an object in Noll's topos) is not injective.

# Engineering challenges

To model topoi on a finite computer, various compromises and trade-offs were needed.

- Although this isn't an iron rule, Bewl's topoi are locally finite, i.e. every  $|\text{Hom}(A, B)| < \infty$
- Digression: It turns out that this condition on a topos implies it has unique injective hulls. This result seems new. I wrote it up as a pure math paper on arXiv
- Bewl also caches products and exponents. For objects  $A$  and  $B$ ,  $A \times B$  and  $B \wedge A$  are computed just once.
- The word "object" is overloaded in computing, so in Bewl, there are "dots" and "arrows".

## Engineering challenges (2)

- I use the 'cake' pattern to define a trait *Topos* as a stack of traits adding helper methods on top of *BaseTopos*
- For example, Bewl can calculate coproducts and coequalizers from the other topos operations.  
This is a verbatim transcript of constructions in McLarty/Moerdijk & Maclane from pure math into software.
- Every dot in Bewl has a type attached to it. So a DOT[T] can be loosely thought of as ranging over values of type T. Functions of type  $T \Rightarrow U$  are easily interchangeable with arrows  $T > U$  (Scala sugar for the Bewl type  $>[T, U]$ ).
- The main difference between functions and arrows is that arrows know their source and target, and can be compared for equality.

## Engineering challenges (3)

- Scala was an almost perfect fit because of its terse style, expressive idioms, and advanced type system.  
(I first tried writing Bewl in Java, then Clojure)

For example, if dotA is a DOT[A] and dotB is a DOT[B], we can construct a new arrow like this:

```
val arrow: A > B =  
  dotA(dotB) { a =>  
    // ...  
    <expression of type B>  
  }
```

We can also apply arrows directly as if they were functions:

```
val a: A = ...  
val b: B = arrow(a)
```

# What are the values over which a dot ranges?

- A topos object (or "dot") in Bewl is a  $\text{DOT}[A]$ , and calculations with it involve manipulating values of type  $A$ , as if the dot somehow ranged over values of type  $A$ . But dots are not sets, and don't have elements.
- These values have meaning only inside the scope of an arrow definition. It's all consistent with the very precise definition of the internal language as described in McLarty's book.
- An earlier version of the DSL interpreted the values  $a: A$  of a  $\text{DOT}[A]$  as arrows  $R \rightarrow A$ , for some "domain of definition" object  $R$ . Now they are pure syntax.

# Tight integration

- Scala supports expressive DSLs
- Some quite involved categorical calculations - for example, the tensorial strength axioms for strong monads - can be described elegantly in Bewl.
- I've considered even deeper integration of the DSL with the language, using Scala implicit magic to make types interchangeable with dots, and functions with arrows.
- In summary, this all works a bit like the (mythical) "category Hask", as a meeting ground of software with math.

# Truth values

- Since the topos may not be Boolean, Bewl has a type called TRUTH which essentially generalizes *Boolean*.
- The 'subobject classifier' in a topos is a DOT[TRUTH].
- Bewl autocalculates the logical operations on TRUTH values (and, or, implies, not) so you can do the equivalent of Boolean algebra.
- In fact, as the subobject classifier of a topos, DOT[TRUTH] is endowed with the structure of a Heyting algebra.

Which brings us to algebraic structures in Bewl.

# Algebraic structures in Bewl

- The DSL lets you define these yourself, using off-the-peg operations (+, ~, etc) with known arities.

Example: Here's the definition of a commutative magma, a structure with one binary operation and one algebraic law:

```
val commutativeMagmas =  
  AlgebraicTheory(*)( $\alpha * \beta := \beta * \alpha$ )  
case class CommutativeMagma[T](  
  override val carrier: DOT[T],  
  op: BinaryOp[T]  
) extends commutativeMagmas.Algebra[T]  
  (carrier)(* := op)
```

- The library includes definitions for monoids, groups, rings, actions, modules, lattices and Heyting algebras.

## Algebraic structures (2)

Definition of a group in Bewl:

```
lazy val groups = AlgebraicTheory(  
  ι, ~, *  
)(  
  "left unit" law ( ι * α := α ),  
  "right unit" law ( α * ι := α ),  
  "left inverse" law ( (~α) * α := ι ),  
  "associative" law (  
    (α * β) * γ := α * (β * γ )  
  )  
)
```

As in the previous example, there's also a case class *Group* to add syntactic sugar.

## Algebraic structures (3)

There are helper functions for building structures:

```
private val (i, x, y) = ('i, 'x, 'y)

val monoidOf3 =
  monoidFromTable(
    i, x, y,
    x, x, y,
    y, x, y
  ) // right-dominant on two generators
```

Bewl can also extract the group of units from a monoid.

## Algebraic structures (4)

- Actions and modules are slightly more involved - they define a new structure within the context of an existing one (monoids / groups and rings, respectively). The DSL caters for this via a concept of 'auxiliary scalars'.
- There are methods to calculate endomorphism monoids and automorphism groups, as algebraic structures in the topos.
- The library could be extended to add many familiar algebraic constructions (e.g. abelianizing a group) which translate naturally in Bewl.

# Topos implementations

Implementing topoi in Bewl is nontrivial.

There are four built-in implementations:

- Finite sets
- The topos of actions of a monoid
- The topos of actions of a group
- The topos of automorphisms

## Topos implementations (2)

The last three all work inside an existing topos.

So if  $\mathcal{E}$  is a topos, we can construct a new topos  $\text{Aut}(\mathcal{E})$  consisting of all the dots-with-a-single-automorphism in  $\mathcal{E}$ .

Similarly, if  $M$  is a monoid object in  $\mathcal{E}$ , we can construct the topos of all objects  $A$  in  $\mathcal{E}$  that come with an action of  $M$ , i.e. an arrow  $A \times M \rightarrow A$  subject to the algebraic laws.

Exponentials, the subobject classifier, logical operations on truth values, etc will all be computed automatically.

# Sample helper methods

Calculating the "name" of an arrow (see McLarty):

```
trait Arrow[S <: ~, T <: ~] ... {  
  ...  
  final lazy val name: UNIT > (S → T) =  
    (source > target).transpose(I) {  
      (i, x) => arrow(x)  
    }  
}
```

From the same class, a method to tell if an arrow is epic:

```
final lazy val isEpic: Boolean =  
  target.exists(source) {  
    (t, s) => target.=?=(  
      t, arrow(s)  
    )  
  } toBool
```

## Sample helper methods (2)

Similarly concise, generic library code uses the DSL to compute:

- is an arrow monic / epic / iso?
- the inverse of an arrow, if it has one
- is an object injective?
- Coterminator (0), coproducts, coequalizers
- the arrow to any object from 0
- epi-mono factorizations

These work in any topos and for me, are a major proof-of-concept validation for Bewl.

# Performance

Better than you'd think, although there are no grounds for complacency. Simple calculations with small sets and monoid actions are fast.

I have largely managed to write code that is clean, efficient and DSL-compliant, thanks to the caching of common operations such as product and exponent.

Most of the required optimizations don't break any abstract layers and have been neatly packed away into 'driver extensions' for specific topoi.

## Performance (2)

Remaining pain points:

- Quantifiers
- Computations with monoid actions

The latter were much improved by a special algorithm which efficiently calculates a presentation for a monoid action (!) so that Bewl can enumerate morphisms between two actions. This is obviously specific to sets.

The algorithm depends on a useful, but very elementary, criterion for finite generating sets of an monoid action.

# Who might use this project?

Bewl definitely needs users. Some possible interest groups:

- People who want to use it as a learning aid to understand category theory. On the GitHub repo, I explain how to quickly set up a command-line REPL for this.
- Music theorists, continuing the Noll approach
- Monoid and semigroup theorists, to explore the possibilities of topos-theoretic reasoning with actions
- anybody who understands both Scala and topoi, and wants to contribute or deepen their understanding!

## Future directions

There is much more to do, but I'd especially like to add

- support for Lawvere-Tierney topologies and sheaves
- construct the topos of coalgebras for a left exact monad
- construct the double-exponential monad for an algebra

These are in principle fairly mechanical, and perhaps the hardest part is to set up decent test fixtures.

Volunteers welcome!

## Future directions (2)

So many promising possibilities, so little time.

<http://github.com/fdilke/bewl>

THANK YOU