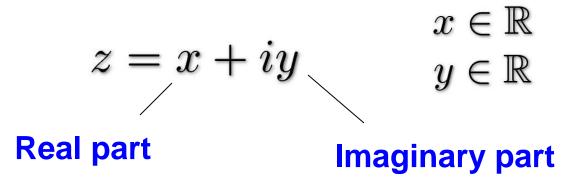
Maths for Computing

Lecture 3

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\mathbb{C} : The complex numbers

• Every complex number can be written as



• We can operate on complex numbers like usual (BODMAS), with the additional rule $i^2 = -1$

BB Dividing complex numbers

 $\frac{a+ib}{c+id} = ?$

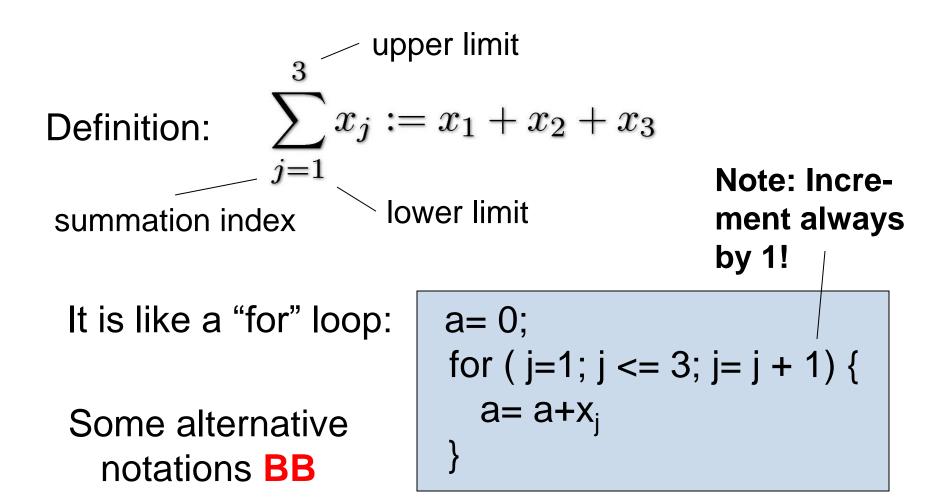
Trick:

$$(c+id) \cdot (c-id) = c^2 - i^2 d^2 - cid + cid = c^2 + d^2$$

Therefore,

$$\begin{aligned} \frac{a+ib}{c+id} &= \frac{a+ib}{c+id} \cdot \frac{c-id}{c-id} = \frac{1}{c^2+d^2}(a+ib)(c-id) \\ &= \frac{1}{c^2+d^2}(ac+bd-iad+icb) \end{aligned}$$

Summation notation (Σ notation)



BB Some alternative notations

$$\sum_{i \in \{1, \dots, N\}} x_i = \sum_{i=1}^N x_i$$

$$\sum_{i \in \{1,...,N\}, i \neq 2} x_i = \left(\sum_{i=1}^N x_i\right) - x_2$$

$$\sum_{\substack{i \in \{1,...,N\}\\j \in \{1,...,M\}}} x_i x_j = \sum_{i=1}^N \sum_{j=1}^M x_i x_j$$

Summation notation

Summation notation is extremely convenient and is used everywhere.

Note: Empty sums are zero:

$$\sum_{k=2}^{1} k^2 = 0$$

Examples: **BB**

BB "Normal" examples

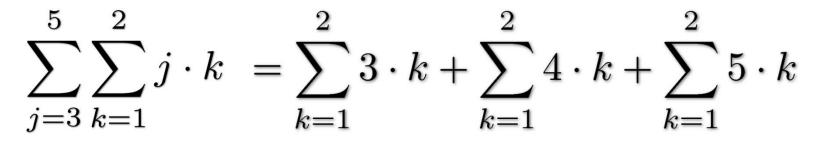
$$\sum_{\substack{i=1\\5\\j=3}}^{4} 3 \cdot i = 3 + 6 + 9 + 12$$

$$\sum_{\substack{j=3\\j=3}}^{5} \sum_{k=1}^{2} j \cdot k = \sum_{\substack{j=3\\j=3}}^{5} (j \cdot 1 + j \cdot 2)$$

$$= 3 \cdot 1 + 3 \cdot 2 + 4 \cdot 1 + 4 \cdot 2 + 5 \cdot 1 + 5 \cdot 2$$

= 3 + 6 + 4 + 8 + 5 + 10

BB More examples



 $= 3 \cdot 1 + 3 \cdot 2 + 4 \cdot 1 + 4 \cdot 2 + 5 \cdot 1 + 5 \cdot 2$

= 3 + 6 + 4 + 8 + 5 + 10

Note how the order of expansion does not matter! Reason:

$$a + (b + c) = (a + b) + c$$
 (Associativity)
 $a + b = b + a$ (Commutativity)

Product notation

Definition:
$$\prod_{j=1}^{5} a_j := a_1 \cdot a_2 \cdot a_3 \cdot a_4 \cdot a_5$$

Think about a for loop, but multiplication inside, instead of summation.

Example:
$$\prod_{j=1}^{5} j = 1 \cdot 2 \cdot 3 \cdot 4 \cdot 5 = 120$$

Product notation

Product notation is somewhat less common but also still used frequently (and useful)

Note: Empty products are 1:

$$\prod_{i=5}^{1} x_i = 1$$

This makes a lot of sense:

Empty sum

If the loop is never executed -> sum is 0.

Empty product

If the loop is never executed -> product is 1.

Why one implies the other: **BB**

BB ... and why one implies the other

For $a, b \in \mathbb{R}$

$$a \cdot b = \exp(\log(a \cdot b)) = \exp(\log(a) + \log(b))$$

$$\prod_{i=1}^{n} x_i = \exp\left(\sum_{i=1}^{n} \log(x_i)\right)$$

If the product is empty, then the sum is empty and

$$\prod_{i=1}^{n} x_i = \exp(0) = 1$$

Summary

Sum: $\sum_{j=1}^{3} x_j := x_1 + x_2 + x_3$, empty sum is 0. Product: $\prod_{j=1}^{5} a_j := a_1 \cdot a_2 \cdot a_3 \cdot a_4 \cdot a_5$, empty product is 1

Multiple sums or products:

$$\sum_{j=3}^{5} \sum_{k=1}^{2} j \cdot k$$

... order of evaluation does not matter

"The domino effect"

INDUCTION

Mathematical induction

Let P(n) be the predicate (or assumption).

P(n) is true for particular values of n, e.g.

P(n) : " $n^3 - n$ is divisible by 3 "

$$n = 1 : 1^{3} - 1 = 0$$

$$n = 2 : 2^{3} - 2 = 8 - 2 = 6$$

$$n = 3 : 3^{3} - 3 = 27 - 3 = 24$$

Proof by induction

$$P(n)_{-}:`n^3-n$$
 is divisible by $\ 3`$

Seems alright for the examples.

But we would like to show: P(n) is true for all $n \in \mathbb{N}$ In "maths speak": $\forall n \in \mathbb{N} : P(n)$ "for all"

Proof by induction

implies

Intuitively ...

- 1. Show that $P(n) \Rightarrow P(n+1)$ for any n
- 2. Show that P(1) is true
- 3. Then from 1. and 2., P(2) is true
- 4. ... and from 1. and 3., P(3) is true
- 5. ... and from 1. and 4., P(4) is true
- 6. ...

Mathematical induction

In "maths speak": "and"

$$P(1) \land (\forall n \in \mathbb{N} : (P(n) \Rightarrow P(n+1)))$$

 $\Rightarrow \forall n \in \mathbb{N} : P(n)$ implies

• If P(1) is true

... and for all n, $P(n) \Rightarrow P(n+1)$, then P(n) is true for all values of n.

• This is the principle of **mathematical induction.**

Example 1

P(n) : ' $n^3 - n$ is divisible by 3'

• Case n=1:
$$1^3 - 1 = 0$$
 ... is divisible by 3

- Now assume that P(n) is true
 - Need to show P(n+1) is true
 - I.e. we need to show that

 $(n+1)^3 - (n+1)$ is divisible by 3. But we can use that $n^3 - n$ is divisible by 3.



Example 2

- Francesco Maurolico (1575)
- First known proof by induction.



$$P(n)$$
 : $\sum_{i=1}^{n} (2i-1) = 1 + 3 + 5 + \ldots + 2n - 1 = n^2$

Example 2 continued

- P(1) : $1 = 1^2$... is true.
- Assume P(n) is true, need to show P(n+1) is then also true.
- Write P(n+1) in terms of P(n):

$$P(n+1)$$
 : $\sum_{i=1}^{n+1} 2i - 1 = \sum_{\substack{i=1 \ =n^2 \text{ by Ind.}}}^{n} 2i - 1 + 2(n+1) - 1$

$$= n^{2} + 2n + 1 = (n + 1)^{2}$$

Example 2 concluded

- So, if P(n) is true, P(n+1) is true
- And P(1) is true.
- Therefore, by mathematical induction, $\Rightarrow \forall n \in \mathbb{N} : P(n)$

BB Example 3

"Geometric series"

Claim:

$$P(n) : \sum_{i=0}^{n} q^{i} = 1 + q + q^{2} + q^{3} + \ldots + q^{n}$$

$$= \frac{1 - q^{n+1}}{1 - q}$$

BB Example 3: Start

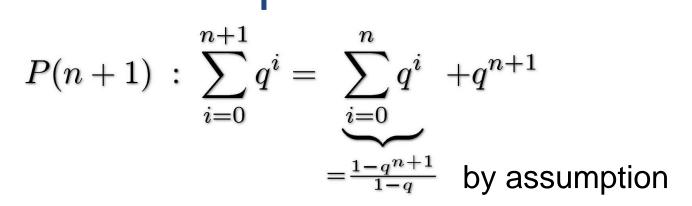
$$P(1) : rac{1-q^2}{1-q} = rac{(1-q)\cdot(1+q)}{1-q}$$

 $= 1+q = \sum_{i=0}^{1} q^i$

 \dots so P(1) is true.

Assume now P(n) is true.

BB Example 3: Induction step



$$= \frac{1 - q^{n+1}}{1 - q} + q^{n+1} \frac{1 - q}{1 - q}$$
$$= \frac{1 - q^{n+1} + q^{n+1} - q^{n+2}}{1 - q}$$
$$= \frac{1 - q^{n+2}}{1 - q} = \frac{1 - q^{(n+1)+1}}{1 - q}$$

BB Example 3: Conclusion

By induction it follows that

 $\sum_{i=0}^n q^i = \frac{1-q^{n+1}}{1-q} \quad \forall \quad n \in \mathbb{N}$

Proof complete. Qed (quod erat demonstrandum)

Things about induction

- Mathematical induction can <u>only</u> be used to prove arguments for positive, whole numbers i.e., the **natural numbers** \mathbb{N} .
- No need to start with P(1). Sometimes it is easier to start with P(3) or P(4) etc.
- But then the inductive proof will only be valid for P(>=3) or P(>=4) etc.

Induction and recursion

- Mathematical induction is important in computer science because it allows us to prove things about recursive programs.
- A recursive function is a function that calls itself.
- Recursion is the basis of several programming languages, e.g., PROLOG.



Recursive programming

- A technique for simplifying a problem by dividing it into subproblems of the same type.
- For example, consider this function of the factorial:

```
function n= Factorial(nm)
if (nm <= 1)
    n= 1;
else
    n= nm*Factorial(nm-1);
end
end</pre>
```

Proof of correctness

Proof by induction:

Induction:

3.

2. Now assume it is correct for n;

else

is correct by assumption = (n-1)!

n= nm*Factorial(nm-1);
end

Proof of correctness complete

Furthermore,

$$\begin{array}{l} n! = n \cdot (n-1) \cdot (n-2) \cdot \ldots \cdot 2 \cdot 1 = n \cdot (n-1)! \\ \\ \dots \\ \\ \text{else} \\ n = n \texttt{m}^* \texttt{Factorial} (\texttt{nm-1}); \\ \\ \text{end} \end{array} \\ \qquad \qquad = (n-1)! \end{array}$$

Therefore, n= nm! ... the function does always return the correct value. Done.

Practical tips about Induction

- If you don't use the assumption, you are not doing proof by induction! BB1
- Sometimes it is useful to "work from two sides" BB2
- If you forget any of the steps it is not a proof!
 BB3
 (consolute the start is often forgetten)

(especially the start is often forgotten)

BB1 Using the assumption

Example 2 again:

We are trying to do the Induction step to show

$$\sum_{i=1}^{n+1} 2i - 1 = (n+1)^2$$

Induction step: $\sum_{i=1}^{n+1} 2i - 1 = \sum_{i=1}^{n} 2i - 1 + \text{something}$ This is "n" -> we can use the assumption here!

BB2 "Coming from two sides"

We need to show

$$\sum_{i=1}^{n+1} 2i - 1 = (n+1)^2$$

"From the left":

$$\sum_{i=1}^{n+1} 2i - 1 = \left(\sum_{i=1}^{n} 2i - 1\right) + 2(n+1) - 1$$

using the assumption simplifying

 $= n^{2} + 2(n+1) - 1 = n^{2} + 2n + 1$

BB "Coming from two sides ..."

"From the right":

$$(n+1)^2 = n^2 + 2n + 1$$

$$\uparrow$$
simplifying

"From the left": $n^2 + 2n + 1$

It's equal – success!

"From the right": $n^2 + 2n + 1$ Wrong

VVrong

BB Forgotten start

"Proof" that for all numbers n = n - 1:

- 1. (forgotten start)
- **2.** Assumption: True for n, i.e. n = n 1
- **3.** Induction:

$$n+1 = (n-1)+1 = n$$
 q.e.d
 \uparrow
Using assumption:
 $n = n - 1$

Wrong

Summary

- Proof by induction always has 3 parts:
- 1. Start: e.g. show P(n) for n=0, or n=1
- **2.** Assumption: P(n) is true
- Induction step: show that P(n+1) is true using P(n)