

SEMANTIC EXPRESSIVE CAPACITY WITH BOUNDED MEMORY

ANTOINE VENANT

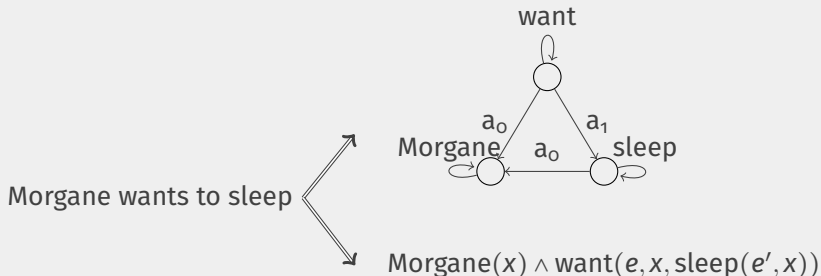
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JOINT WORK WITH ALEXANDER KOLLER

SEMANTIC INTERPRETATION

Linguistic expression \Rightarrow (formal) meaning representation.
Representations can be logical formulae, or graphs (AMR [Banarescu & all 2013], MRS [Copestake & all 2005]).



- Consensual approach: semantic interpretation is a *compositional* process, guided by syntax.

Statement

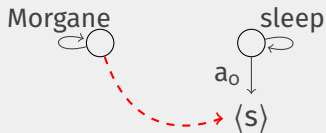
“The meaning of a complex expression is a **function** of the meaning of its parts and the **syntactic rule** that combines them.”

Requires:

- A syntax tree, along which semantic construction is performed in a bottom-up fashion.
- Operators for semantic composition (semantic algebra).
- Which semantic interpretation functions can we express *compositionally* using specific classes of syntax trees **and** semantic operators?

(Essentially) one job:

combine predicates with their arguments.



TWO TRADITIONS

'Unification style'

Finite set of markers denoting 'holes' ($\langle s \rangle$, $\langle o \rangle$, $\langle \text{mod} \rangle$, $\langle \text{comp} \rangle$) waiting to be filled with semantic values. Markers accessible in unconstrained order [Copestake & all, 2001, Courcelles & Englefriet 2012, Groshwitz & all 2017].

'Lambda style'

Countably infinite ordered set of markers but order constrain access (variables' scope) [Montague 1977, Steedman 2001].

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→ number of 'holes' accessible at a given time of the construction process is bounded: '**bounded memory**'.

'Lambda style'

Countably infinite ordered set of markers but order constrain access (variables' scope) [Montague 1977, Steedman 2001].

'bounded memory' operators are popular

- In semantic parsing [Chiang & all 2013, Groschwitz & all 2018, Chen & all 2018].
- For the manual design of grammars [Bender 2002 *inter alia*].

Expressive limitation due to bounded memory capacity?

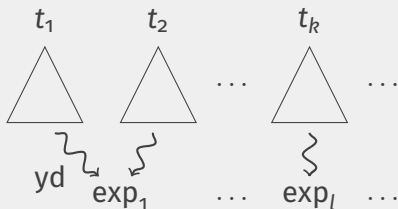
- Specifically, considering long distance dependencies.
- If impossible (from distance) to combine a predicate with its argument right away → need to store argument slot until argument becomes available.

FURTHER MOTIVATION

- A lot is known on expressive capacity of grammatical formalisms – in terms of languages (of words/trees).
 - *e.g.*, famous CCG/TAG/LIG [Vijay-Shanker & Weir, 1994] weak equivalence result.
- What about the joint expressivity of grammatical formalisms and specific semantic combinators in terms of *relations*?
- **Do (weakly) equivalent grammatical formalisms support the same compositional interpretations?**
- Inform the elaboration of semantic parsing systems

ABSTRACT VIEW ON GRAMMARS

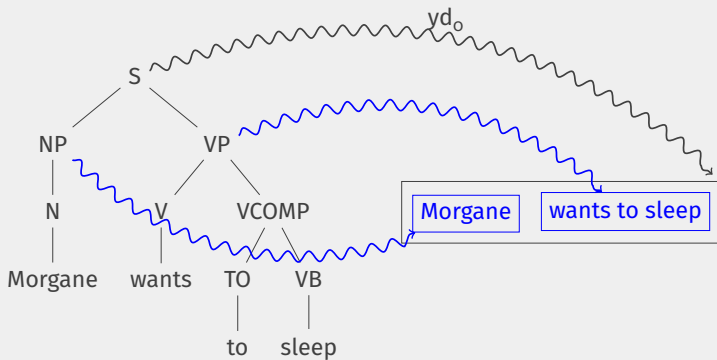
- Set of 'grammatical' syntax trees $\{t_1, t_2, \dots\}$.
- *yield* function, yd , associating each tree with its string projection (the linguistic expression for which it is a grammatical analysis).



- The set $\{t_1, t_2, \dots\}$ could be given by any kind of descriptive/computing device (formal grammar, neural net, ...).

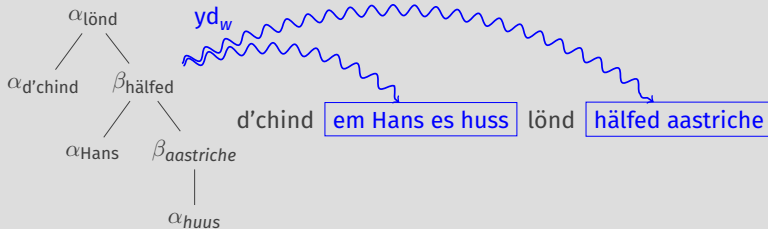
THE PROJECTIVE YIELD yd_0

- Concatenates children's yield from left to right.



A NON-PROJECTIVE YIELD: yd_w

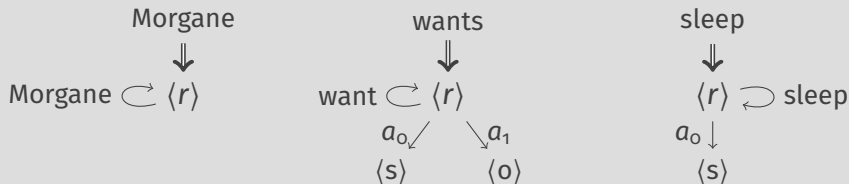
Swiss-German cross-serial dependencies [Shieber 1985]



(dass) (mer) d' chind em Hans es huus lönd
(that) (we) the-children-ACC Hans-DAT the-house-ACC let
hälfed aastriche
help paint

'(that we) let the children help Hans paint the house'

Interpretation for elementary syntactic constituents



- $\langle s \rangle$, $\langle o \rangle$, $\langle r \rangle$: markers.
- $\langle s \rangle$, $\langle o \rangle$: argument placeholders ('holes'): a semantic value will eventually be substituted for them during the process of semantic composition.
- $\langle r \rangle$: root of the semantic constituent ('hook'), destined to be substituted for an argument placeholder.

Semantic algebra (i.e. composition operators)

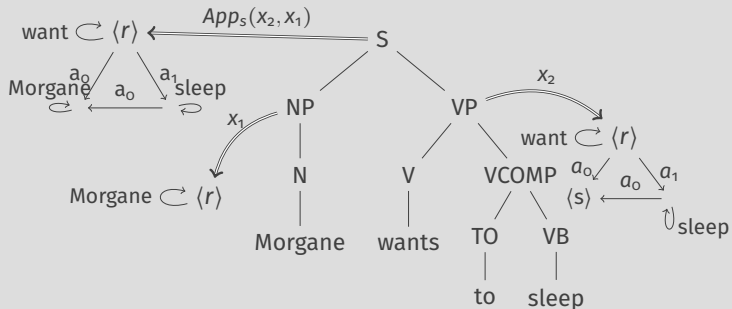
- Example with the AM algebra [Groschwitz & all 2017]

$$App_o \left(\begin{array}{c} \text{want} \curvearrowright \langle r \rangle \\ a_o \swarrow \quad \searrow a_1 \\ \langle s \rangle \quad \langle o \rangle \end{array} \mid \begin{array}{c} \langle r \rangle \curvearrowright \text{sleep} \\ a_o \downarrow \\ \langle s \rangle \end{array} \right) = \begin{array}{c} \text{want} \curvearrowright \langle r \rangle \\ a_o \swarrow \quad \searrow a_1 \\ \langle s \rangle \quad \langle s \rangle \end{array} \begin{array}{c} \curvearrowright \text{sleep} \\ \downarrow \\ \text{sleep} \end{array}$$

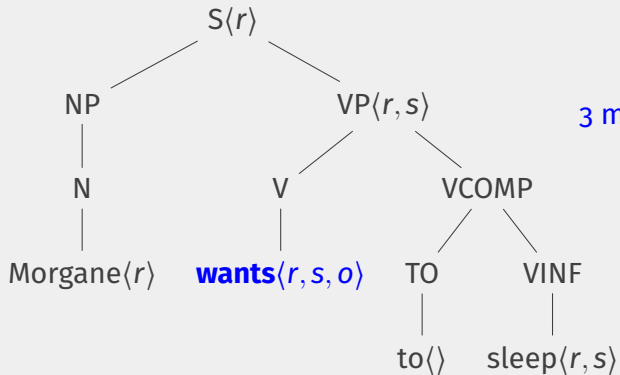
The diagram illustrates the application of the operator App_o to two semantic structures. The left structure shows a root $\langle r \rangle$ (want) and a root $\langle r \rangle$ (sleep) connected by a dashed blue arrow. A marker $\langle o \rangle$ is shown in red, connected to the root $\langle r \rangle$ (want) by a dashed red arrow. The right structure shows the result of the application, where the marker $\langle o \rangle$ has been merged with the root $\langle r \rangle$ (sleep), and the root $\langle s \rangle$ has been merged with the root $\langle s \rangle$.

- Merge referenced marker $\langle o \rangle$ of the functor with the root $\langle r \rangle$ of the argument, then ‘forgets’ these two markers.
- Merge any other identical marker (here, $\langle s \rangle$).

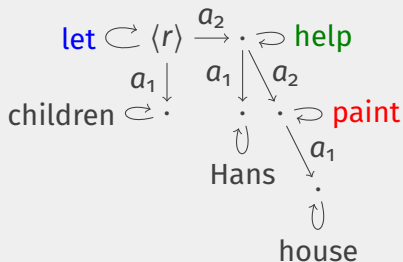
Homomorphic interpretation of syntax trees

$$\{VP(x_1, x_2) \rightarrow APP_o(x_1, x_2), S(x_1, x_2) \rightarrow APP_s(x_2, x_1)\}$$


'SEMANTIC' MEMORY



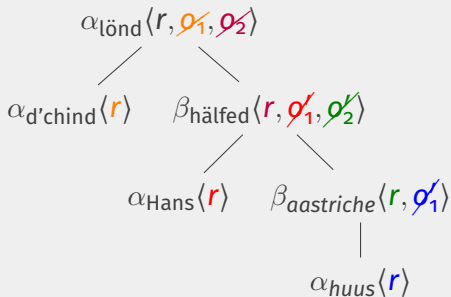
3 markers required



d'chind em Hans es huss lönd hälfed aastriche

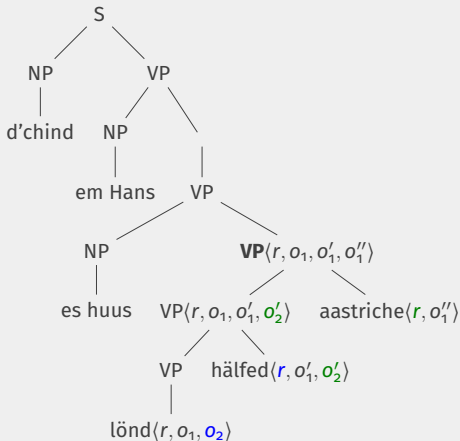
- **lönd**: $\langle r, o_1, o_2 \rangle$ (two objects).
- **hälfed**: $\langle r, o_1, o_2 \rangle$ (two objects).
- **aastriche**: $\langle r, o_1 \rangle$ (one object).

Non-projective analysis possible with a 3-markers capacity.



PROJECTIVITÉ AND MEMORY 3/3

With a projective analysis: 4 markers seem intuitively required.



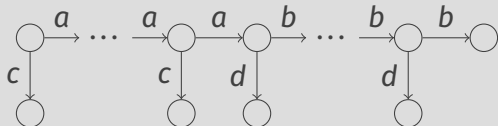
ABSTRACTING AWAY

- Arbitrary long crossed-serial dependencies \rightarrow infinite memory required?
- A formal relation for a mathematical proof:

CSD

Word to graph function $w \mapsto g_w$ where

- Words w are of the form: $\underbrace{a \dots a}_n \underbrace{b \dots b}_m \underbrace{c \dots c}_n \underbrace{d \dots d}_m$.
n times m times n times m times
- And for each such w , g_w is:



Theorem ?

There exists no projective grammar and finite memory compositional interpretation mechanism over a projective grammar which expresses CSD.

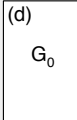
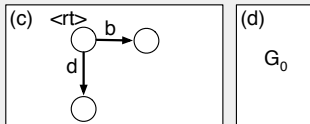
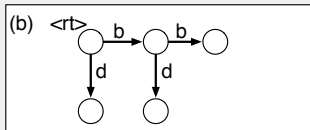
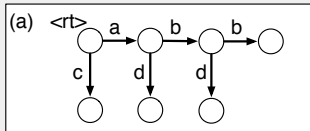
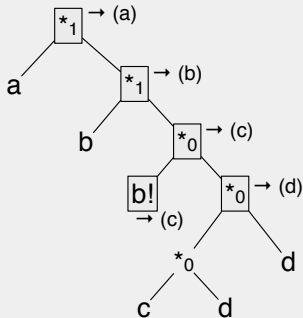
NOT A Theorem

There exists no projective grammar and finite memory compositional interpretation mechanism over a projective grammar which expresses CSD.

UNNATURAL CONSTRUCTIONS

NOT A Theorem

There exists no projective grammar and finite memory compositional interpretation mechanism over a projective grammar which expresses CSD.



THEOREM

If one further impose specific alignments between elementary syntactic and semantic constituents (' a ' aligned with ' \underline{a} ', ' b ' aligned with ' \underline{b} ' ...) it can be shown:

Theorem

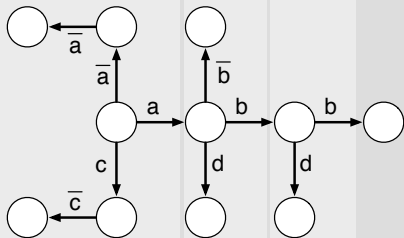
- **there exists no projective grammar and finite memory compositional interpretation mechanism over a projective grammar** expressing CSD and respecting elementary alignments.
- There exists a **non-projective** grammar and a finite-memory compositional interpretation mechanism expressing CSD and respecting elementary alignments.
- Remark: strong assumption on alignments but no assumption on grammatical formalism.

IMPERFECT ALIGNMENTS

- Without the alignment condition the theorem is false.
- However, weaker form of alignments can be achieved if we constrain the grammatical formalism (pumping lemma).
- **Requires arbitrary complex 'arguments' to avoid previous unnatural constructions.**
- Result for Tree-Adjoining Grammars (TAG).

$\overline{\text{CSD}}$ relation

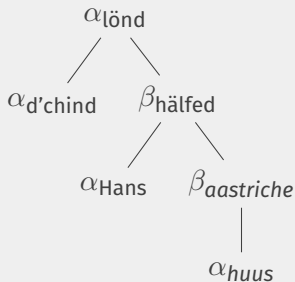
$a\overline{a}b\overline{b}c\overline{c}d\overline{d}$



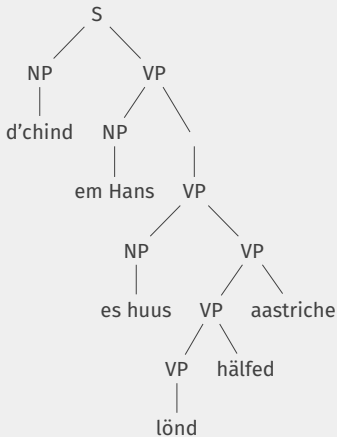
TWO KINDS OF TREES

- TAG grammars produce derivation trees and derived trees.

Derivation (dependency) tree

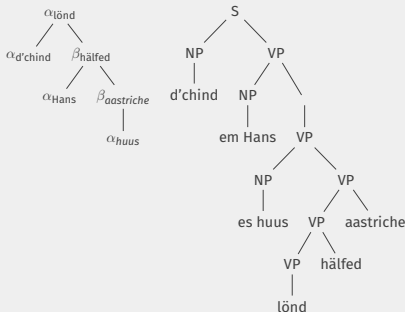


Derived (syntagmatic) tree



TWO (WEAKLY) EQUIVALENT GRAMMAR FORMALISMS

- Formalism TAG: Use the **derivation trees** of some TAG grammar with a **non-projective** yield.
- Formalism PTAG: Use the **derived trees** of some TAG grammar with the **projective** yield.
- The two formalisms generate the same **word languages**, but not necessarily the same **relations**.



Theorem

- There exists a (non-projective) TAG grammar and a finite memory compositional interpretation mechanism expressing $\overline{\text{CSD}}$.
- There exists no (projective) PTAG grammar and finite memory compositional interpretation mechanism expressing $\overline{\text{CSD}}$.

- Theoretical result on the link between compositionality, projectivity and bounded memory capacity.
- Strong result, under strong assumption of perfect syntax/semantics alignments at the lexical level.
- **independent** of considered grammatical formalism.
- New light shed on the choice between derivation/derived tree as the support of semantic composition for TAG grammars.
- Do weakly equivalent grammatical formalisms support the same compositional interpretation mechanisms? → **No!**

CONCLUSIONS AND FUTURE WORK

- Notion of expressivity at the syntax/semantics interface.
- Theoretical study on the link between projectivity and 'semantic' memory.
- What could we say about more restricted forms of non-projectivity? Finite increase in required memory capacity?
- Artificial non-projectivity due to imperfect aligners in semantic parsing systems.
- Locally translate from 'unification style' to 'lambda style' to circumvent projectivity issues?

Questions?